

# CSE 140: Components and Design Techniques for Digital Systems

## Lecture 6: Universal Gates

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# Combinational Logic: Various Types of Gates

- Universal Set of Gates
  - Motivation
  - Definition
  - Examples
- Other Types of Gates
  - XOR
  - NAND / NOR
  - Block Diagram Transfers

# Universal Set of Gates: Motivation

- AND, OR, NOT: Logic gates related to reasoning from Aristotle (384-322BCE).
- NAND, NOR: Inverted AND, Inverted OR gates. For VLSI technologies, all gates are inverted (AND,OR operation with a bubble at output).
- XOR: Exclusive OR gates. Parity check.
- Multiplexer + input table: Table based logic for programmability. FPGA technology.
- Neuron and Synapse: Neural network.
- Reversible Gates: Quantum computing.

In the future, we may have new sets of gates due to new technologies. Given a set of gates, can the gates in the set cover all possible switching functions?

# Universal Set

Universal set is a powerful concept to identify the coverage of a set of gates afforded by a given technology.

Criterion: If the set of gates can implement AND, OR, and NOT gates, the set is universal.

# Universal Set Definition

Universal Set: A set of gates such that every switching function can be implemented with the gates in this set.

Examples

{AND, OR, NOT}

{AND, NOT}

{OR, NOT}

# Universal Set: Examples

Universal Set: A set of gates such that every switching function can be implemented with the gates in this set.

Examples

{AND, OR, NOT}

{AND, NOT} OR can be implemented with AND &

NOT gates:  $a + b = (a'b')'$

{OR, NOT} AND can be implemented with OR &

NOT gates:  $ab = (a' + b')'$

{XOR} is not universal

{XOR, AND} is universal

# iClicker

Is the set {AND, OR} (but no NOT gate) universal?

A. Yes

B. No

Note that the set was used in a once popular design style as **domino logic** for high performance computing.

# iClicker

Is the set  $\{f(x, y) = xy'\}$  universal?

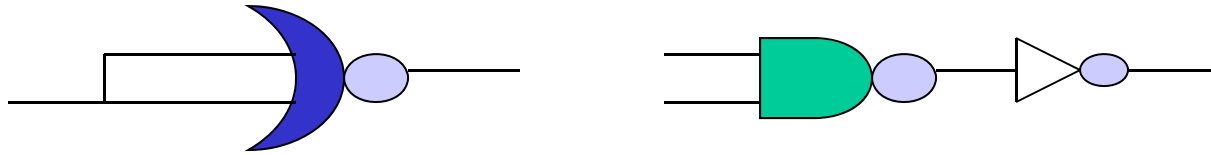
A. Yes

B. No

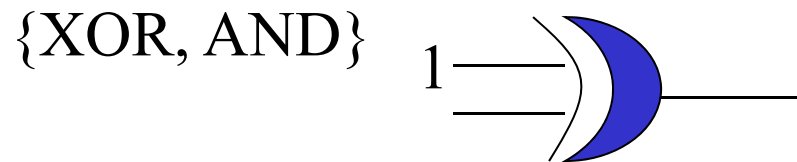


# Universal Set: Examples

{NAND, NOR}



{XOR}



$$x \oplus 1 = x1' + x'1 = x' \text{ if constant "1" is available.}$$

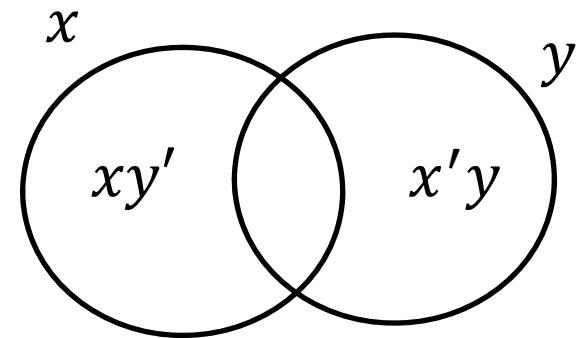
## Other Types of Gates: Properties and Usage

1. XOR  $x \oplus y = xy' + x'y$
2. NAND, NOR
3. Block Diagram Transfers

## Other Types of Gates: XOR

$$x \oplus y = xy' + x'y$$

It is a parity function (examples)  
useful for testing because the  
flipping of a single input changes the  
output.

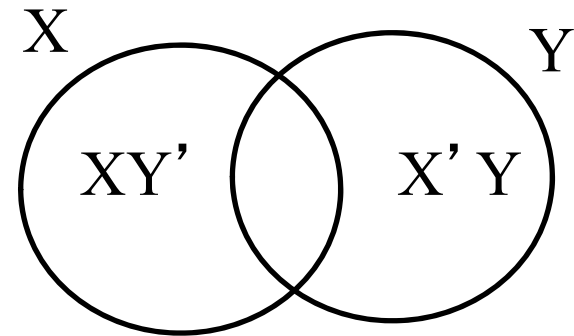


id	x	y	$x \oplus y$
0	0	0	0
1	0	1	1
2	1	0	1
3	1	1	0

	x=0	x=1
y=0	0	1
y=1	1	0

## Other Types of Gates: XOR

1) XOR  $x \oplus y = xy' + x'y$

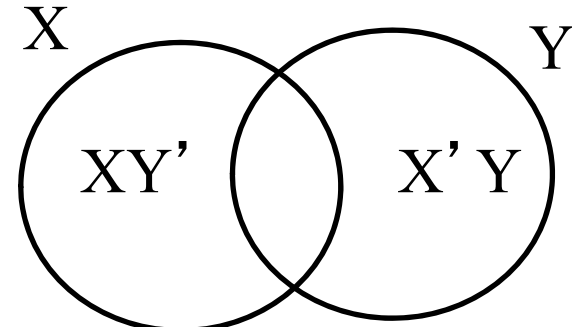


(a) Commutative  $x \oplus y = y \oplus x$

(b) Associative  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$

## Other Types of Gates: XOR

1) XOR  $x \oplus y = xy' + x'y$



a) Commutative  $x \oplus y = y \oplus x$

b) Associative  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$

c)  $1 \oplus x = x', 0 \oplus x = x$

d)  $x \oplus x = 0, x \oplus x' = 1$

## Other Types of Gates: Properties and Usage

e) If  $ab = 0$ , then  $a \oplus b = a + b$

Proof: If  $ab = 0$ , then

$$a = a(b + b') = ab + ab' = ab'$$

$$b = b(a + a') = ba + ba' = a'b$$

$$\text{Therefore, } a + b = ab' + a'b = a \oplus b$$

Note that in full adder, we have

$$c_{\text{out}} = ab + bc + ac = ab + c(a + b)$$

From property e), we can also write  $c_{\text{out}} = ab + c(a \oplus b)$

## Other Types of Gates: XOR

f)  $f(x, y) = x \oplus xy' \oplus x'y \oplus (x + y) \oplus x = ?$

(Priority of operations: AND,  $\oplus$ , OR)

Hint: We apply Shannon's Expansion.

# Shannon's Expansion (for switching functions)

**Formula:**  $f(x, y) = xf(1, y) + x'f(0, y)$

Proof by enumeration:

If  $x = 1$ ,  $f(x, y) = f(1, y)$ :

$$xf(1, y) + x'f(0, y) = 1f(1, y) + 1'f(0, y) = f(1, y)$$

If  $x = 0$ ,  $f(x, y) = f(0, y)$ :

$$xf(1, y) + x'f(0, y) = 0f(x, y) + 0'f(0, y) = f(0, y)$$



## Other types of gates: XOR

Simplify the function (Priority of operations: AND,  $\oplus$ , OR)

$$f(X,Y) = X \oplus XY' \oplus X'Y \oplus (X+Y) \oplus X$$

$$\text{Case } X = 1: f(1, Y) = 1 \oplus Y' \oplus 0 \oplus 1 \oplus 1 = Y$$

$$\text{Case } X = 0: f(0, Y) = 0 \oplus 0 \oplus Y \oplus Y \oplus 0 = 0$$

Thus, using Shannon's expansion, we have

$$f(X, Y) = Xf(1,Y) + X'f(0,Y) = XY$$

# XOR gates

iClicker: Is the equation

$$a+(b\oplus c) = (a+b)\oplus(a+c) \text{ true?}$$

A. Yes

B. No

## Other Types of Gates: NAND, NOR

2) NAND, NOR gates

NAND (NOR) gates are not associative

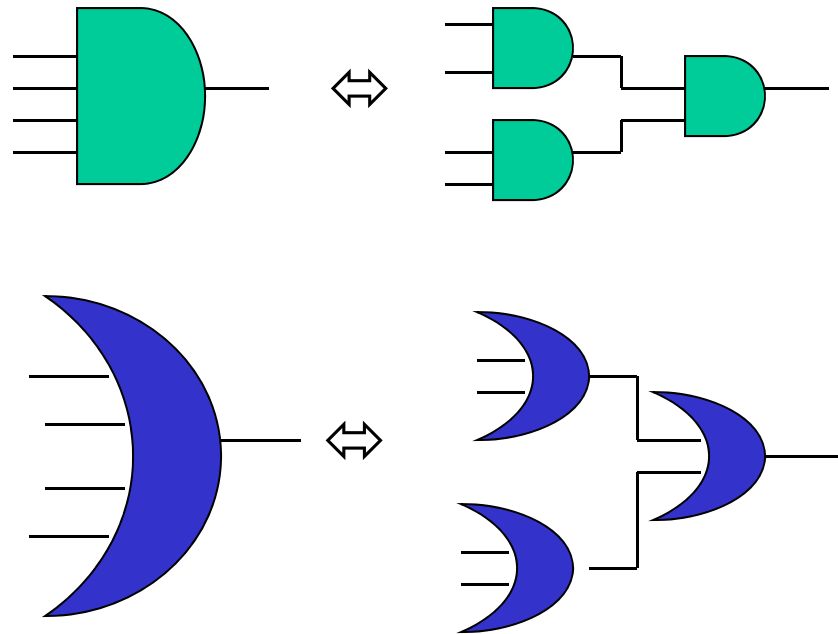
Let  $a \mid b = (ab)'$

$$(a \mid b) \mid c \neq a \mid (b \mid c)$$

# Other Types of Gates: Block Diagram Transform

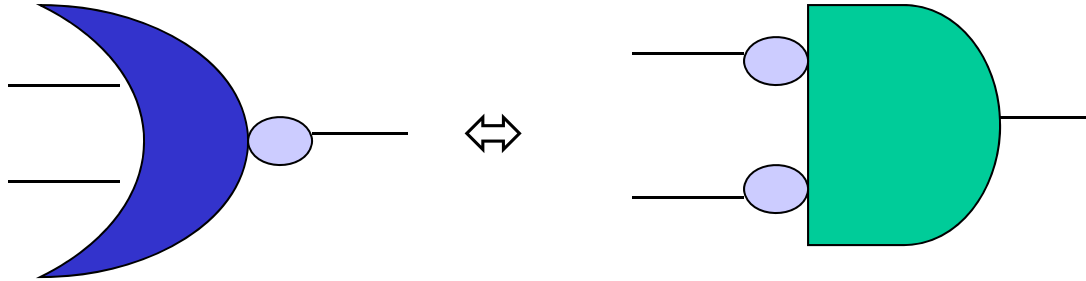
## 3) Block Diagram Transformation

a) Reduce # of inputs.

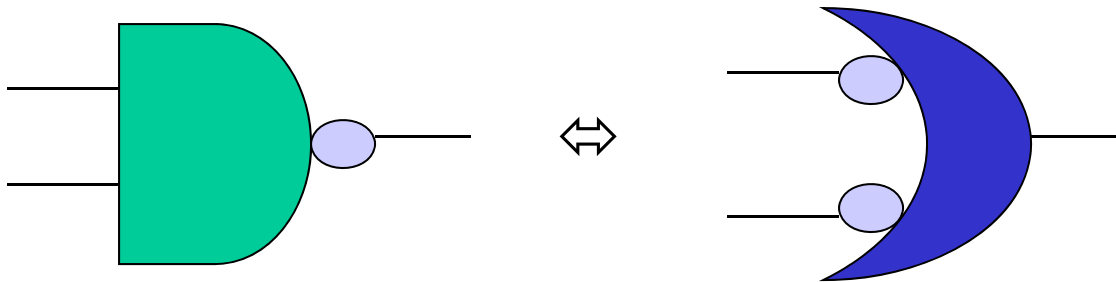


# Other Types of Gates: Block Diagram Transform

## b. DeMorgan's Law



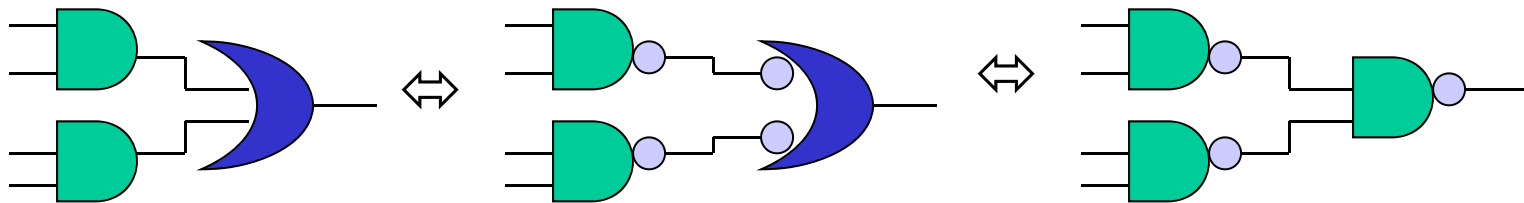
$$(a+b)' = a' b'$$



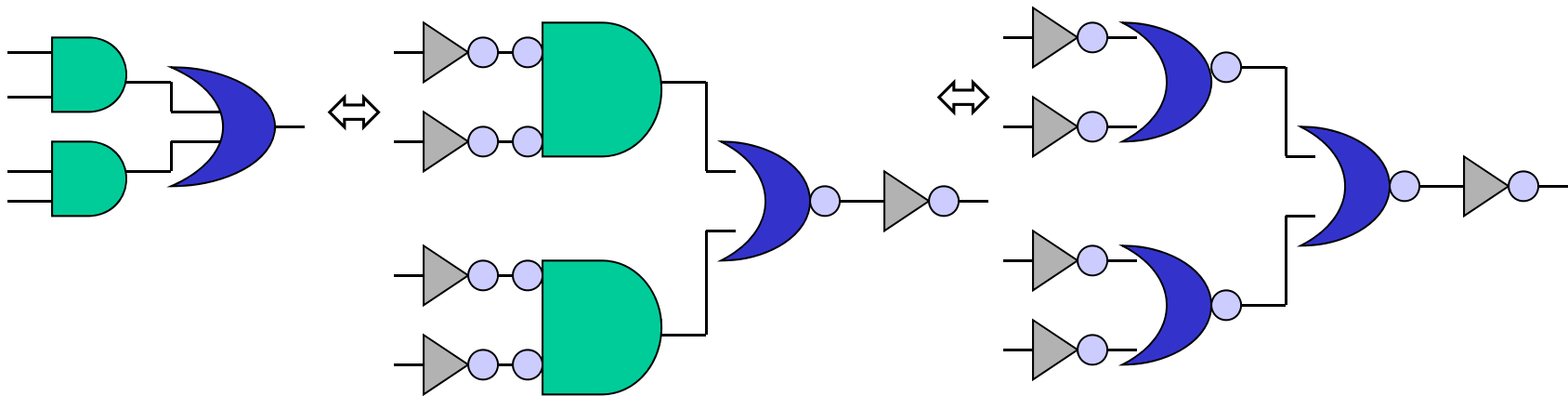
$$(ab)' = a' + b'$$

# Other Types of Gates: Block Diagram Transform

## c. Sum of Products (Using only NAND gates)

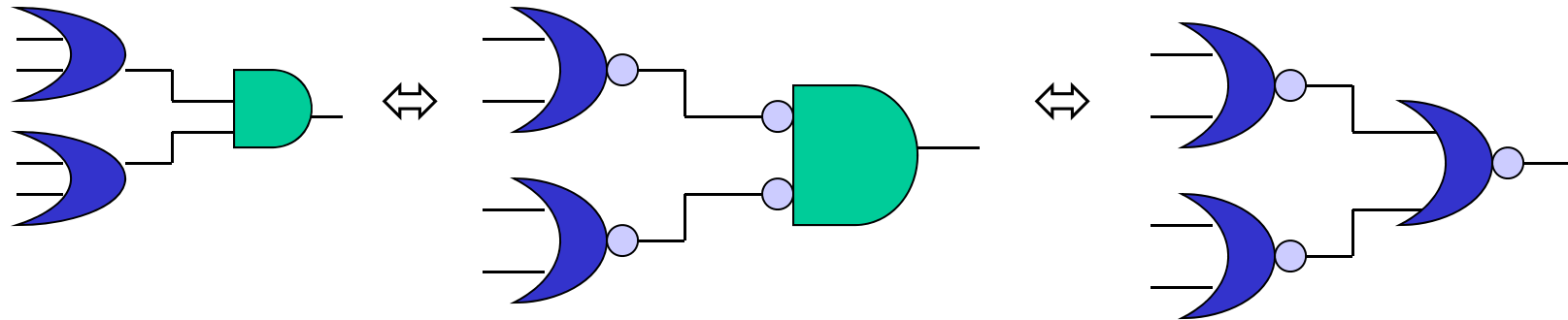


## Sum of Products (We create many bubbles with NOR gates)



# Other Types of Gates: Block Diagram Transform

d. Product of Sums (NOR gates only)



We will create many bubbles with NAND gates.

# Other Types of Gates: Block Diagram Transform

## NAND, NOR gates

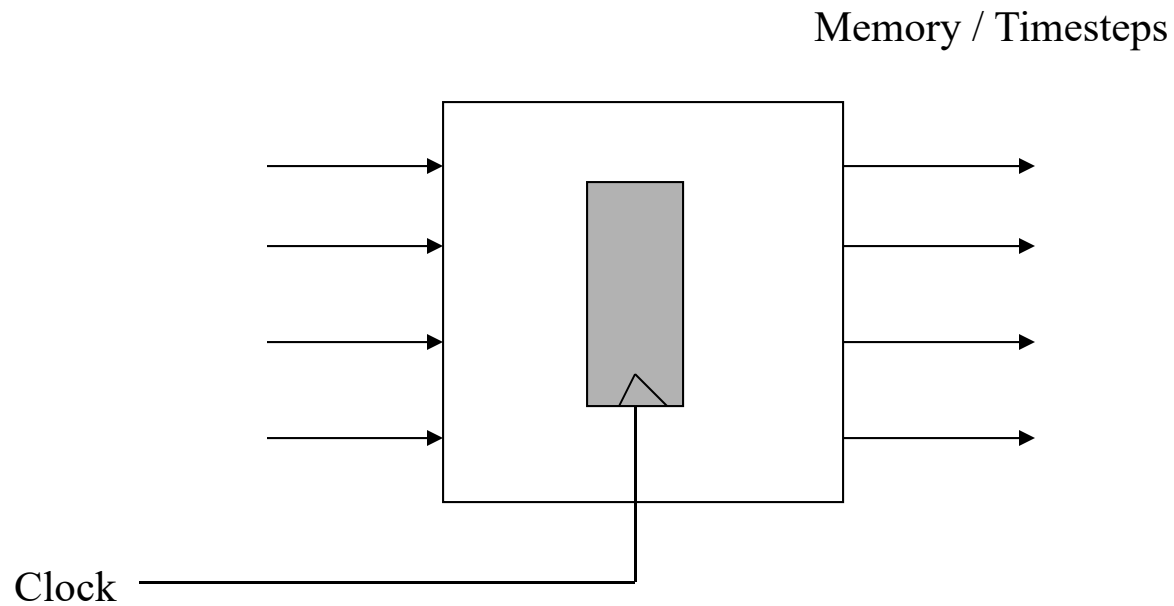
Remark:

Two level NAND gates: Sum of Products

Two level NOR gates: Product of Sums



## Part II. Sequential Networks



Flip flops  
Specification  
Implementation